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Advanced Functional Programming

02 - Testing

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• When is a program correct?

- When is a program correct?
- What is a specification?
- How to establish a relation between the specification and the implementation?
- What about bugs in the specification?

- "Equals can be substituted for equals"
- In other words: if an expression has a value in a context, we can replace it with any other expression that has the same value in the context without affecting the meaning of the program.
- When we deal with infinite structures: two things are equivalent if we cannot find out about their difference:

- "Equals can be substituted for equals"
- In other words: if an expression has a value in a context, we can replace it with any other expression that has the same value in the context without affecting the meaning of the program.
- When we deal with infinite structures: two things are equivalent if we cannot find out about their difference:

ones = 1: ones ones' = 1:1: ones' In most functional languages like ML or OCaml, there is no referential transparency:

let val x = ref 0
fun f n = (x := !x + n; !x)
in f 1 + f 2

In most functional languages like ML or OCaml, there is no referential transparency:

let val x = ref 0
fun f n = (x := !x + n; !x)
in f 1 + f 2

But we cannot replace the last line with 1 + f 2, even though f 1 = 1.

• Haskell is referentially transparent – all side-effects are tracked by the IO monad.

do

```
x <- newIORef 0
let f n = do modifyIORef x (+n); readIORef x
r <- f 1
s <- f 2
return (r + s)</pre>
```

Note that the type of f is Int \rightarrow IO Int – we cannot safely make the substitution we proposed previously.

Because we can safely replace equals for equals, we can *reason* about our programs – this is something you already saw in the course on functional programming.

For example to prove some statement P xs holds for all lists xs, we need to show:

- P [] the base case;
- for all x and xs, P xs implies P (x:xs).

- Equational reasoning can be an elegant way to prove properties of a program.
- Equational reasoning can be used to establish a relation between an "obviously correct" Haskell program (a specification) and an efficient Haskell program.
- Equational reasoning can become quite long...
- Careful with special cases (laziness):
 - undefined values;
 - partial functions;
 - infinite values.

You can formalize such proofs in other systems such as Agda, Coq or Isabelle.

QuickCheck, an automated testing library/tool for Haskell

Features:

- Describe properties as Haskell programs using an embedded domain-specific language (EDSL).
- Automatic datatype-driven random test case generation.
- Extensible, e.g. test case generators can be adapted.

- Developed in 2000 by Koen Claessen and John Hughes.
- Copied to other programming languages: Common Lisp, Scheme, Erlang, Python, Ruby, SML, Clean, Java, Scala, F#
- Erlang version is sold by a company, QuviQ, founded by the authors of QuickCheck.

```
isort :: Ord a \Rightarrow [a] \rightarrow [a]
isort [] = []
isort (x:xs) = insert x (isort xs)
```

```
insert :: Ord a \Rightarrow a \rightarrow [a] \rightarrow [a]

insert x [] = [x]

insert x (y:ys)

| x \leq y = x : y : ys

| otherwise = y : insert x ys
```

We can now try to prove that for all lists xs,

length (sort xs) = length xs.

- The base case is trivial.
- The inductive case requires a lemma relating insert and length suggestions?

Consider the following (buggy) implementation of insertion sort:

```
sort :: [Int] \rightarrow [Int]
sort [] = []
sort (x:xs) = insert x xs
```

```
insert :: Int \rightarrow [Int] \rightarrow [Int]
insert x [] = [x]
insert x (y:ys) | x \leq y = x : ys
| otherwise = y : insert x ys
```

Let's try to debug it using QuickCheck.

A good specification is

- as precise as necessary,
- no more precise than necessary.

A good specification for a particular problem, such as sorting, should distinguish sorting from all other operations on lists, without forcing us to use a particular sorting algorithm.

Certainly, sorting a list should not change its length.

```
sortPreservesLength :: [Int] → Bool
sortPreservesLength xs =
  length (sort xs) = length xs
```

We can test by invoking the function :

```
> quickCheck sortPreservesLength
Failed! Falsifiable, after 4 tests:
[0,3]
```

```
sort :: [Int] \rightarrow [Int]

sort [] = []

sort (x:xs) = insert x xs

insert :: Int \rightarrow [Int] \rightarrow [Int]
```

insert x [] = [x] insert x (y:ys) | x ≤ y = x : ys | otherwise = y : insert x ys

Which branch does not preserve the list length?

> quickCheck sortPreservesLength
OK, passed 100 tests.

Looks better. But have we tested enough?

Properties are first-class objects

```
(f `preserves` p) x = p x = p (f x)
```

sortPreservesLength = sort `preserves` length

idPreservesLength = id `preserves` length

Properties are first-class objects

```
(f `preserves` p) x = p x = p (f x)
```

sortPreservesLength = sort `preserves` length

idPreservesLength = id `preserves` length

So id also preserves the lists length:

> quickCheck idPreservesLength
OK, passed 100 tests.

We need to refine our spec.

We can define a predicate that checks if a list is sorted:

isSorted	:: [Int] -	\rightarrow	Bool	
isSorted	[]	=	True	
isSorted	[×]	=	True	
isSorted	(x:y:xs)	=	x < y & isSorted (y:xs)	

And use this to check that sorting a list produces a list that isSorted.

```
> quickCheck sortEnsuresSorted
Falsifiable, after 5 tests:
[5,0,-2]
> sort [5,0,-2]
[0,-2,5]
```

We're still not quite there...

What's wrong now?

```
sort :: [Int] \rightarrow [Int]
sort [] = []
sort (x:xs) = insert x xs
```

insert :: Int \rightarrow [Int] \rightarrow [Int]

What's wrong now?

```
sort :: [Int] \rightarrow [Int]
sort [] = []
sort (x:xs) = insert x xs
```

insert :: Int \rightarrow [Int] \rightarrow [Int]

We are not recursively sorting the tail in sort.

Another bug

```
> quickCheck sortEnsuresSorted
Falsifiable, after 7 tests:
[4,2,2]
```

```
> sort [4,2,2]
[2,2,4]
```

This is correct. What is wrong?

Another bug

```
> quickCheck sortEnsuresSorted
Falsifiable, after 7 tests:
[4,2,2]
```

```
> sort [4,2,2]
[2,2,4]
```

This is correct. What is wrong?

```
> isSorted [2,2,4]
False
```

The isSorted spec reads:

sorted :: [Int] \rightarrow Bool sorted [] = True sorted (x:[]) = True sorted (x:y:ys) = x < y & sorted (y : ys)

Why does it return False? How can we fix it?

Are we done yet?

Is sorting specified completely by saying that

- sorting preserves the length of the input list,
- the resulting list is sorted?

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Is sorting specified completely by saying that

- · sorting preserves the length of the input list,
- the resulting list is sorted?

No, not quite.

```
evilNoSort :: [Int] → [Int]
evilNoSort xs = replicate (length xs) 1
```

This function fulfills both specifications, but still does not sort.

We need to make the relation between the input and output lists precise: both should contain the same elements – or one should be a permutation of the other.

```
permutes :: ([Int] \rightarrow [Int]) \rightarrow [Int] \rightarrow Bool
permutes f xs = f xs `elem` permutations xs
```

```
sortPermutes :: [Int] \rightarrow Bool
sortPermutes xs = sort `permutes` xs
```

This completely specifies sorting and our algorithm passes the corresponding tests.

To use QuickCheck in your program:

import Test.QuickCheck

Define properties.

Then call to test the properties.

```
quickCheck :: Testable prop \Rightarrow prop \rightarrow IO ()
```

The type of is an *overloaded* type:

```
quickCheck :: Testable prop \Rightarrow prop \rightarrow IO ()
```

- The argument of is a property of type prop
- The only restriction on the type is that it is in the Testable type class.
- When executed, prints the results of the test to the screen hence the result type.

So far, all our properties have been of type :

```
sortPreservesLength :: [Int] \rightarrow Bool
sortEnsuresSorted :: [Int] \rightarrow Bool
sortPermutes :: [Int] \rightarrow Bool
```

When used on such properties, QuickCheck generates random integer lists and verifies that the result is True.

If the result is for 100 cases, this success is reported in a message.

If the result is False for a test case, the input triggering the result is printed.

Other example properties

```
appendLength :: [Int] \rightarrow [Int] \rightarrow Bool
appendLength xs ys =
length xs + length ys = length (xs ++ ys)
```

```
plusIsCommutative :: Int \rightarrow Int \rightarrow Bool
plusIsCommutative m n = m + n = n + m
```

```
takeDrop :: Int \rightarrow [Int] \rightarrow Bool
takeDrop n xs = take n xs ++ drop n xs = xs
```

```
dropTwice :: Int \rightarrow Int \rightarrow [Int] \rightarrow Bool
dropTwice m n xs =
```

drop m (drop n xs) = drop (m + n) xs

> quickCheck takeDrop

```
OK, passed 100 tests.
```

> quickCheck dropTwice
Falsifiable after 7 tests.
1
-1
[0]

```
> drop (-1) [0]
[0]
```

```
> drop 1 (drop (-1) [0])
[]
```

Nullary properties

A property without arguments is also possible:

```
lengthEmpty :: Bool
lengthEmpty = length [] = 0
```

wrong :: Bool
wrong = False

> quickCheck lengthEmpty
OK, passed 100 tests.

> quickCheck wrong
Falsifiable, after 0 tests.

No random test cases are involved for nullary properties.

QuickCheck subsumes unit tests.

Recall the type of quickCheck:

```
quickCheck :: Testable prop \Rightarrow prop \rightarrow IO ()
```

We can now say more about when types are Testable:

• testable properties usually are functions (with any number of arguments) resulting in a Bool

What argument types are admissible?

QuickCheck has to know how to produce random test cases of such types.

```
class Testable prop where
  property :: prop → Property
instance Testable Bool where
  ...
instance (Arbitrary a, Show a, Testable b) ⇒
    Testable (a → b) where
```

We can test any Boolean value or any testable function for which we can generate arbitrary input.

```
collect :: (Testable prop, Show a) \Rightarrow
a \rightarrow prop \rightarrow Property
```

The function gathers statistics about test cases. This information is displayed when a test passes:

```
> let sPL = sortPreservesLength
> quickCheck (\xs -> collect (null xs) (sPL xs))
OK, passed 100 tests.
96% False
```

4% True.

The result implies that not all test cases are distinct.

26% 0.

- 21% 1.
- 15% 2.

10% 5.

10% 3.

• • •

Most lists are small in size: QuickCheck generates small test cases first, and increases the test case size for later tests.

In the extreme case, we can show the actual data that is tested:

```
> quickCheck (\ xs → collect xs (sPL xs))
OK, passed 100 tests:
6% []
1% [9,4,-6,7]
1% [9,-1,0,-22,25,32,32,0,9,...
...
```

Why is it important to have access to the test data?

The function insert preserves an ordered list:

```
implies :: Bool \rightarrow Bool \rightarrow Bool
implies x y = not x || y
```

insertPreservesOrdered :: Int → [Int] → Bool insertPreservesOrdered x xs = sorted xs `implies` sorted (insert x xs)

Implications – contd.

> quickCheck insertPreservesOrdered OK, passed 100 tests.

But:

OK, passed 100 tests.

88% False

12% True

For 88 test cases, insert has not actually been relevant.

Implications – contd.

The solution is to use the QuickCheck implication operator:

```
(\Longrightarrow) :: (Testable prop) \Rightarrow
Bool \rightarrow prop \rightarrow Property
```

instance Testable Property

The type allows to write a logically equivalent formula that also explicitly rejects the test case.

```
iPO :: Int \rightarrow [Int] \rightarrow Property
iPO x xs = sorted xs \implies sorted (insert x xs)
```

Now, lists that are not sorted are discarded and do not contribute towards the goal of 100 test cases.

Implications – contd.

We can now easily run into a new problem:

```
iPO :: Int \rightarrow [Int] \rightarrow Property
iPO x xs = length xs > 2 & sorted xs \implies
sorted (insert x xs)
```

We try to ensure that lists are not too short, but:

The chance that a random list is sorted is extremely small. QuickCheck will give up after a while if too few test cases pass the precondition.

Configuring QuickCheck

```
quickCheckWith :: Testable prop \Rightarrow Args \rightarrow prop \rightarrow IO ()
```

```
data Args where
  replay :: Maybe (StdGen, Int)
  -- should we replay a previous test?
  maxSuccess :: Int
  -- max number of successful tests
  -- before succeeding
  maxDiscardRatio :: Int
  -- max number of discarded tests
  -- per successful test
  maxSize :: Int
  --max test case size
```

Generators

- Instead of increasing the number of test cases to generate, it is usually better to write a custom random generator.
- Generators belong to an abstract data type Gen. Think of as a restricted version of IO. The only effect available to us is access to random numbers.
- We can define our own generators using another domain-specific language. The default generators for datatypes are specified by defining instances of class Arbitrary:

```
class Arbitrary a where
  arbitrary :: Gen a
  ...
```

choose	::	Random $a \Rightarrow (a,a) \rightarrow Gen a$
oneof	::	[Gen a] → Gen a
frequency	::	$[(Int, Gen a)] \rightarrow Gen a$
elements	::	[a] → <mark>Gen</mark> a
sized	::	(Int \rightarrow Gen a) \rightarrow Gen a

Simple generators

```
instance Arbitrary Bool where
  arbitrarv = choose (False, True)
instance (Arbitrary a, Arbitrary b) \Rightarrow
            Arbitrary (a,b) where
  arbitrary = do
                  x <- arbitrarv
                  v <- arbitrary</pre>
                  return (x,y)
```

data Dir = North | East | South | West
instance Arbitrary Dir where

arbitrary = elements [North, East, South, West]

• A simple possibility:

instance Arbitrary Int where

arbitrary = choose (-20, 20)

• Better:

instance Arbitrary Int where

arbitrary = sized (\ n \rightarrow choose (-n,n))

• QuickCheck automatically increases the size gradually, up to the configured maximum value.

Idea: Adapt the default generator for lists.

The following function turns a list of integers into a sorted list of integers:

For example:

> mkSorted [1,2,-3,4]
[1,3,6,10]

The generator can be adapted as follows:

There is another function to construct properties provided by QuickCheck, passing an explicit generator:

```
forAll :: (Show a, Testable b) \Rightarrow
Gen a \rightarrow (a \rightarrow b) \rightarrow Property
```

This is how we use it:

```
iPO :: Int \rightarrow Property
iPO x = forAll genSorted
(\ xs \rightarrow length xs > 2 & sorted xs \implies
sorted (insert x xs))
```

Arbitrary revisited

```
class Arbitrary a where
arbitrary :: Gen a
shrink :: a \rightarrow [a]
```

The other method in is

```
shrink :: (Arbitrary a) \Rightarrow a \rightarrow [a]
```

- Maps each value to a number of 'structurally smaller' values.
- When a failing test case is discovered, is applied repeatedly until no smaller failing test case can be obtained.

Program coverage

To assess the quality of your test suite, it can be very useful to use GHC's program coverage tool:

```
$ ghc -fhpc Suite.hs --make
$ ./Suite
$ hpc report Suite --exclude=Main --exclude=QC
   18% expressions used (30/158)
    0% boolean coverage (0/3)
         0\% guards (0/3), 3 unevaluated
       100% 'if' conditions (0/0)
       100% gualifiers (0/0)
        . . .
```

This also generates a .html file showing which code has (not) been executed.

modulo	Top Level Definitions			Alternatives		Expressions	
module	%	covered / total	%	covered / total	%	covered / total	
module <u>Prettify2</u>	42%	9/21	23%	8/34	18%	30/158	
Program Coverage Total	42%	9/21	23%	8/34	18%	30/158	

Figure 1: screenshot

```
data Doc = Empty
           Char Char
           Text String
           Line
           Concat Doc Doc
         | Union Doc Doc
         deriving (Show, Eq)
{-- /snippet Doc --}
instance Monoid Doc where
    mempty = empty
    mappend = (<>)
{-- snippet append --}
empty :: Doc
(<>) :: Doc -> Doc -> Doc
{-- /snippet append --}
empty = Empty
Empty <> y = y
x \iff Empty = x
x \ll y = \hat{x} Concat y
char :: Char -> Doc
char c = Char c
```

Figure 2: screenshot

- Haskell can deal with infinite values, and so can QuickCheck. However, properties must not inspect infinitely many values. For instance, we cannot compare two infinite values for equality and still expect tests to terminate. Solution: Only inspect finite parts.
- QuickCheck can generate functional values automatically, but this requires defining an instance of another class Coarbitrary but showing functional values is problematic.
- QuickCheck has facilities for testing properties that involve I0, but this is more difficult than testing pure properties.

QuickCheck is a great tool:

- A domain-specific language for writing properties.
- Test data is generated automatically and randomly.
- Another domain-specific language to write custom generators.
- Use it!

However, keep in mind that writing good tests still requires training, and that tests can have bugs, too.

Required:

- Chapter 11 of Real World Haskell
- *QuickCheck: A Lightweight Tool for Random Testing of Haskell Programs,* Koen Claessen and John Hughes

Additional reading:

- Software Testing with QuickCheck, John Hughes
- Smallcheck and lazy smallcheck: automatic exhaustive testing for small values, Colin Runciman, Matthew Naylor, Fredrik Lindblad
- Hedgehog, Jacob Stanley