



Universiteit Utrecht

[Faculty of Science  
Information and Computing Sciences]

# Talen en Compilers

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# 14. Compiler optimizations



# Announcements

- ▶ Grades soon
- ▶ Next lecture: summary
- ▶ Send topics, questions, example exercises to Lawrence



# This lecture

## Compiler optimizations

Optimization passes

Simple optimizations

Loop optimizations

Other optimizations



# 14.1 Optimization passes



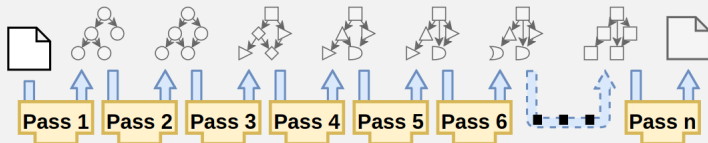
# What is a compiler optimization

- ▶ A terrible name
- ▶ Semantics-preserving code transformation
- ▶ Hopefully improving the code by some metric



# Visualizing optimization passes

Typically the same input and output type, like pass 2.



## 14.2 Simple optimizations





# Peephole optimizations

- ▶ Group of simple but effective optimizations
- ▶ Find and replace
- ▶ Usually on low-level instructions



# Peephole optimizations

- ▶ Group of simple but effective optimizations
- ▶ Find and replace
- ▶ Usually on low-level instructions
  
- ▶  $x * 2 \Rightarrow x \ll 1$
- ▶  $x * 0 \Rightarrow 0$
- ▶  $x \leftarrow 3; x \leftarrow 4 \Rightarrow x \leftarrow 4$



# Unreachable/dead code elimination

- ▶ Uncalled methods/functions
- ▶ Code after a return statement
- ▶ Patterns that cannot be matched
- ▶ ...



# Tail call elimination

```
int add (int m, int n) {  
    if (m = 0) then  
        return n;  
    else  
        return add (m - 1, n + 1);  
}
```

```
int add (int m, int n) {  
    while (m != 0) {  
        m = m - 1;  
        n = n + 1;  
    }  
    return n;  
}
```



## 14.3 Loop optimizations



# Loop optimizations

- ▶ Loop unrolling
- ▶ Loop invariant code motion
- ▶ Loop fusion
- ▶ Loop fission



# Loop unrolling

```
for (int i = 0; i < n; i++)  
{ doStuff (i);  
}
```

```
for (int i = 0; i < n - 4; i += 4)  
{ doStuff (i);  
  doStuff (i + 1);  
  doStuff (i + 2);  
  doStuff (i + 3);  
}
```



# Loop unrolling

```
for (int i = 0; i < n; i++)  
{ doStuff (i);  
}
```

```
for (int i = 0; i < n - 4; i += 4)  
{ doStuff (i);  
  doStuff (i + 1);  
  doStuff (i + 2);  
  doStuff (i + 3);  
}
```

If  $n$  is not divisible by 4, you need to do extra iterations before or after the loop.





# Loop invariant code motion

```
for (int i = 0; i < n; i++)  
{ x = 10 * y + cos (0.5);  
  doStuff (i, x);  
}
```

```
x = 10 * y + cos (0.5);  
for (int i = 0; i < n; i++)  
{ doStuff (i, x);  
}
```



# Loop fusion

```
for (int i = 0; i < n; i++)  
{ doStuff1 (i);  
}  
for (int i = 0; i < n; i++)  
{ doStuff2 (i);  
}
```

```
for (int i = 0; i < n; i++)  
{ doStuff1 (i);  
  doStuff2 (i);  
}
```



# Loop fission

```
for (int i = 0; i < n; i++)  
{ doStuff1 (i);  
  doStuff2 (i);  
}
```

```
for (int i = 0; i < n; i++)  
{ doStuff1 (i);  
}  
for (int i = 0; i < n; i++)  
{ doStuff2 (i);  
}
```



# Loop fission

```
for (int i = 0; i < n; i++)  
{ doStuff1 (i);  
  doStuff2 (i);  
}
```

```
for (int i = 0; i < n; i++)  
{ doStuff1 (i);  
}  
for (int i = 0; i < n; i++)  
{ doStuff2 (i);  
}
```

The opposite of fusion: Sometimes one is better, sometimes the other!

To choose, we might want to add analyses to our compiler.



## 14.4 Other optimizations



# Inlining

```
let x = 5  
in x * y + x
```

```
5 * y + 5
```



# Common Subexpression Elimination

|  $\cos(5x) / (1 + \cos(5x))$  | **let**  $y = \cos(5x)$   
| **in**  $y / (1 + y)$



# Common Subexpression Elimination

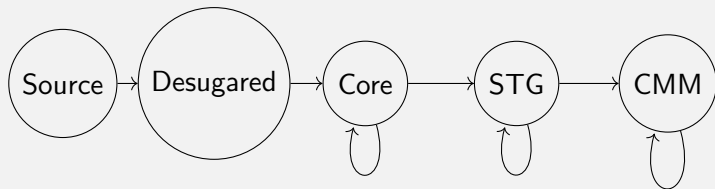
$\cos(5x) / (1 + \cos(5x))$	<b>let</b> $y = \cos(5x)$ <b>in</b> $y / (1 + y)$
-----------------------------	--

Opposite of inlining: Tradeoff between computation and memory.





# Compiler pipeline



## Analysis and code transformation



Analysis and code transformation for optimal fusion



# My research

Analysis and code transformation for optimal fusion of array operations

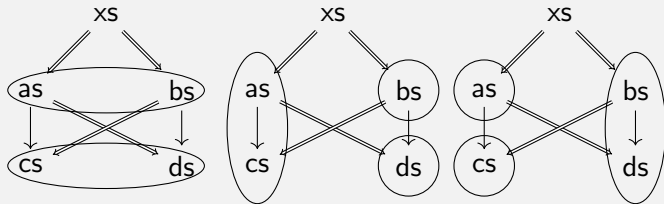


# My research

Analysis and code transformation for optimal fusion of array operations in a data parallel functional language



Analysis and code transformation for optimal fusion of array operations in a data parallel functional language



# Vertical fusion

```
for (int i = 0; i < n; i++) {  
  y [i] = 2 * x [i];  
}  
for (int i = 0; i < n; i++) {  
  z [i] = 4 + y [i];  
}  
return z;
```

```
for (int i = 0; i < n; i++) {  
  z [i] = 4 + 2 * x [i];  
}  
return z;
```

Replacing an array with a scalar

Eliminating  $n$  array reads and writes

